

In the Claims:

Please amend the claims as follows.

1. (Currently Amended) ~~A scalable method for implementing FFT/IFFT computations in multiprocessor architectures that provides improved throughput by eliminating the need for inter-processor communication after the computation of the first " $\log_2 P$ " stages of the FFT/IFFT computations for a multiprocessor architecture including an implementation using " P " processing elements~~A method for controlling processing elements in a multiprocessor architecture to provide improved throughput for FFT/IFFT computations, the method comprising the steps of:

~~Computing~~computing, on a multiprocessor architecture including " P " processing elements, each butterfly of the first " $\log_2 P$ " stages of an FFT/IFFT on either a single one of the processing elements or on each of the " P " processing elements simultaneously; and

distributing the computations of the butterflies in all the subsequent stages of the FFT/IFFT among the " P " processing elements~~processors~~such that each chain of cascaded butterflies consisting of those butterflies that have inputs and outputs connected together, are processed by the same processor~~processing element~~to thereby eliminate the need for inter-processor communication among the processing elements after the computation of the first " $\log_2 P$ " stages of the FFT/IFFT.

2. (Currently Amended) A method as claimed in claim 1 wherein the distributing of the computation of the butterflies subsequent to the first " $\log_2 P$ " butterflies is achieved by assigning operand addresses of each set of butterfly operands to each ~~processor~~processing element in such a manner that the butterfly is processed by the same ~~processor~~processing element that computed the connected butterfly of the previous stage in the same chain of butterflies.

3. (Original) A method as claimed in claim 2 wherein the desired assignment of operand addresses is achieved by deriving the address of the first operand in the operand pair corresponding to the " i^{th} " stage of the computation from the address of the corresponding operand in the previous stage by inserting a "0" in the " $(i+1)^{\text{th}}$ " bit position of the address, while the address of the second operand is derived by inserting a "1" in the " $(i+1)^{\text{th}}$ " bit position of the operand address.

4. (Currently Amended) A method as claimed in claim 1 further including the computing of twiddle factors for the butterfly computations at each ~~processor~~processing element by initializing a counter and then incrementing it by a value corresponding to the number of ~~processors~~processing elements "P" and appending the result with a specified number of "0"s.

5. (Currently Amended) ~~A system for obtaining scalable implementation of FFT/IFFT computations in multiprocessor architectures that provides improved throughput by eliminating the need for inter-processor communication after the computation of the first " $\log_2 P$ " stages of the FFT/IFFT computations for a multiprocessor architecture including an implementation using "P" processing elements~~A system for controlling processing elements in a multiprocessor architecture including P processing elements, comprising:

a means for computing each butterfly of the first " $\log_2 P$ " " $\log_2 P$ " stages of an FFT/IFFT on either a single processing element processor or on each of the "P" processing elements processors simultaneously;

an addressing means for distributing the computation of the butterflies in all the subsequent stages of the FFT/IFFT among the "P" ~~processors~~ processing elements such that each chain of cascaded butterflies consisting of those butterflies that have inputs and outputs connected together, are processed by the same ~~processor~~ processing element to thereby eliminate the need for inter-processor

communication among the processing elements after the computation of the first " $\log_2 P$ " stages of the FFT/IFFT;

means for counting; and

means for computing twiddle factors for the butterfly computations at each ~~processor~~processing elements, the means for computing initializing the means for counting and then incrementing the means for counting by a value corresponding to the number of ~~processors~~processing elements "P" and appending the result with a specified number of "0"s.

6. (Currently Amended) A system as claimed in claim 5 wherein the addressing means comprises addresses generation means for deriving the operand addresses of the butterflies subsequent to the first " $\log_2 P$ " butterflies in such a manner that the butterfly is processed by the same ~~processor~~processing element that computed the connected butterfly of the previous stage in the same chain of butterflies.

7. (Original) A system as claimed in claim 6 wherein the address generation means is a computing mechanism for deriving the address of the first operand in the operand pair corresponding to the " i^{th} " stage of the computation from the address of the corresponding operand in the previous stage by inserting a "0" in the " $(i+1)^{\text{th}}$ " bit position of the address, and deriving the address of the second operand by inserting a "1" in the " $(i+1)^{\text{th}}$ " bit position of the operand address.

8. (Currently Amended) A system as claimed in claim 5 further including a computing mechanism for address generation of twiddle factors for each butterfly on the corresponding ~~processor~~processing element.

9. (Currently Amended) A method of performing a fast Fourier transform or inverse fast Fourier transform on an input signal~~a plurality of inputs to generate a plurality of outputs, the method being performed on a plurality of processors and each transform including a plurality of stages containing at least one butterfly computational block, the method comprising:~~

storing samples of the input signal in a memory;

retrieving the samples from the memory and from these retrieved samples calculating the butterfly computational blocks for the first $\log_2(P)$ ~~$\log_2 P$~~ stages of the transform on a single ~~one of the~~ processors or on a plurality of the processors operating in parallel; and

eliminating the need for communication among and between the processors after the computation of the first " $\log_2 P$ " stages of the transform by calculating chains of butterfly computational blocks corresponding to the subsequent stages of the transform within each of the processors, each chain of butterfly computational blocks that is calculated in a respective processor having inputs and outputs coupled in series.

10. (Currently Amended) The method of claim 9 wherein the first $\log_2(P)$ ~~$\log_2 P$~~ stages of the transform are calculated on all of the processors operating in parallel.

11. (Original) The method of claim 9 wherein the method is performed on two processors, and wherein the first two stages of a radix-2 fast Fourier transform or inverse fast Fourier transform are calculated as a single radix-4 stage, and wherein the subsequent stages of the transform are computed as radix-2 stages.

12. (Original) The method of claim 11 wherein chains comprises a single loop that iterates $N/2 * (\log_2(N/2)) / (\text{number of processors})$ times.

13. (Original) The method of claim 12 wherein each butterfly computational block includes a plurality of operands each having an associated address, and wherein calculating chains of butterfly computational blocks corresponding to the subsequent stages comprises assigning addresses to each of the operands so that each butterfly block in a chain is calculated in the same processor.

14. (Original) The method of claim 13 wherein each butterfly computational block includes a pair of operands, and wherein the operand addresses of these operands are assigned by deriving the address of the first operand in the operand pair corresponding to the " i^{th} " stage of the calculation in the chain from the address of the corresponding operand in the previous stage by inserting a "0" in the " $(i+1)^{\text{th}}$ " bit position of the operand address, and deriving the operand address of the second operand by inserting a "1" in the " $(i+1)^{\text{th}}$ " bit position of the operand address.

15. (Original) The method of claim of claim 9 further comprising initializing a counter and then incrementing the counter by a value corresponding to the number of processors and appending the result with a specified number of "0"s to compute the twiddle factors for each butterfly computational block.

16. (Currently Amended) A processor system, comprising:
a memory operable to store samples of an input signal;
a plurality of processors coupled to the memory, the plurality of processors operable to receive the samples from the memory and to use the samples to execute a fast Fourier transform or inverse fast Fourier transform algorithm on a plurality of inputs to generate a plurality of outputs, each transform including a plurality of stages containing at least one butterfly computational block, and the processors operable to execute the butterfly computational blocks for the first " $\log_2 P$ " " $\log_2 P$ " stages of the a fast Fourier transform or inverse fast Fourier transform on either a

single one of the processors or on a plurality of the processors operating in parallel;
and

address circuitry coupled to the memory and processor and operable to
distribute the computation of the butterfly computational blocks in all stages
subsequent to the first ~~$\log_2 P$~~ $\log_2 P$ states/stages among the plurality of processors
such that each chain of cascaded butterfly computational blocks in the transform are
coupled in series and are computed by the same processor to thereby eliminate the
need for communication among and between the processors after the computation of
the first " $\log_2 P$ " stages of the transform.

17. (Original) The processor system of claim 16 wherein the address circuitry is further operable to derive operand addresses for each of the butterfly blocks subsequent to the first " $\log_2 P$ " butterfly blocks so that each of the butterfly computational blocks is computed by the same processor that computed a butterfly computational block of the previous stage in the same chain of butterfly computational blocks.

18. (Original) The processor system of claim 17 wherein each butterfly computational block includes a pair of operands, and wherein the address circuitry assigns operand addresses of these operands by deriving the address of the first operand in the operand pair corresponding to the " i^{th} " stage of the calculation in the chain from the address of the corresponding operand in the previous stage by inserting a "0" in the " $(i+1)^{\text{th}}$ " bit position of the operand address, and deriving the operand address of the second operand by inserting a "1" in the " $(i+1)^{\text{th}}$ " bit position of the operand address.

19. (Original) The processor system of claim 17 wherein the processors further comprise a counter that is initialized and then incremented by a

value corresponding to the number of processors, an output of the counter being appended with a specified number of "0"s to compute twiddle factors for each butterfly computational block.

20. (Original) The processor system of claim 16 wherein each of the processors comprises a digital signal processor.

21. (Currently Amended) An electronic system, comprising:
a processor system, including,

a memory operable to store samples of an input signal;

a plurality of processors coupled to the memory, the plurality of processors operable to receive samples from the memory and to use the samples to execute a fast Fourier transform or inverse fast Fourier transform algorithm on a plurality of inputs to generate a plurality of outputs, each transform including a plurality of stages containing at least one butterfly computational block, and the processors operable to execute the butterfly computational blocks for the first " $\log_2 P$ " " $\log_2 P$ " stages of the a fast Fourier or inverse fast Fourier transform on either a single one of the processors or on a plurality of the processors operating in parallel; and

address circuitry coupled to the memory and processor and operable to distribute the computation of the butterfly computational blocks in all stages subsequent to the first $\log_2 P$ $\log_2 P$ states stages among the plurality of processors such that each chain of cascaded butterfly computational blocks in the transform are coupled in series and are computed by the same processor to thereby eliminate the need for communication among and between the processors after the computation of the first " $\log_2 P$ " stages of the transform.

22. (Currently Amended) The electronic system of claim 22-21 wherein the electronic system comprises a communications system.

23. (Currently Amended) The electronic system of claim ~~20~~21 wherein each of the processors comprises a digital signal processor.